



## An Improved Variable Length Lightweight DESSHA for IoT-Enabled Smart Devices

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### ABSTRACT

Amidst the accelerating scenarios of cybersecurity and increasing concerns regarding data integrity organizations seriously seek strong and scalable frameworks as safeguard against sensitive data to keep it unmodified from illegitimate access. To ensure consistent performance and optimization in resource utilization over various types of devices security and scalability are of utmost requirements. Data integrity has become challenging issue in resource constrained devices as used in IoT. IoT has become one the major components in automation of healthcare, smart home and smart agriculture. This study proposes a DES based Scalable Hashing Algorithm (DESSHA) for smart devices using IoT. DESSHA due to its close alignment with theoretical values of all security test including sensitivity, confusion diffusion, uniform distribution and collision resistance with scalability. It provides message authentication for nearly all types of smart devices on IoT, especially for smart agriculture and E-Commerce purposes. Our analysis demonstrates that DESSHA provides a high level of security for message authentication for IoT enabled devices.

**Keywords:-**DESSHA, Cryptography, Hashing, SHA 256, Smart devices, Security of IoT.

### 1. INTRODUCTION

Data security has become increasingly important as computer technology has become more widely used. The use of the internet and mobile devices has increased greatly, particularly in studies in areas such as Smart devices on IoT, e-commerce, banking, finance, security, and education [1]. The advancement of technology like software defined networks [2] and software programs utilized today has made it possible for hostile consumers to intercept information at various stages of data transmission[3]. As a result, it is far more critical than ever for any company or individual to protect their sensitive and important documents. To maintain a superior line of defense, security constantly looks for how to keep information safe from unauthorized intrusion[4]. Encryption is a means of preserving data such that it remains unmodified and secure at some point throughout the transfer of information from the sender to the intended recipient [5]. Many encryption technologies have been developed to assure data security and protection [6]. To properly develop and execute security, we must be one step forward, or consider along the same lines as, cyber criminals [7]. Maintaining secrecy against



outsiders has always been a problem, and encryption can assist to keep exchanges of confidential messages hidden. Cryptography encompasses a wide range of procedures, such as microdots, combining words with visuals, and further inventive methods of preventing information from being captured by other parties [8]. Cryptography is used to exchange and process information safely and effectively, without a third-party interception. Encryption, hashing, and steganography are examples of such procedures [9].

Cryptographic hash functions are a method for ensuring the integrity of a document. A cryptographic hash function algorithm compresses a message to create a fingerprint for data integrity purposes [10]. Hash functions are versatile cryptographic building blocks used to ensure data validity and digital signatures. Their main fields of application are password protection, message authentication, digital signatures, certificates, and cryptocurrencies like Bitcoin [11].

The Secure Hash Algorithm (SHA), a family of cryptographic hash functions, was released in 1993 under the title Secure Hash Standard. SHA-1 input can accept up to 64 bits and output can be up to 160 bits long. Because the number of input constraints minimizes collisions that may occur when using these algorithms, this has become one of the algorithm's strengths. SHA-1 processes input in 512-bit blocks and then divides it into 16 subblock's of 32 bits each [12]. Because of advancements in the field of cryptography, particularly the introduction of the AES encryption standard, which offers better security with three key sizes of 128, 192, and 256 bits, the security of SHA-1 is no longer considered adequate [13]. In response, the NSA launched a project to create new hash functions with key sizes of 128, 192, and 256 bits that would provide comparable security to AES. This project resulted in the creation of the Federal Information Processing Standard hash family, which includes three new hash functions: SHA-256, SHA-384, and SHA-512 [14]. SHA-256 algorithm has two steps: preprocessing and hash computation. Padding a message and parsing the padded message into m-blocks are examples of pre-processing. The initialization values are configured to be used in the hash computation. The size of the hash value affects the security of the SHA-256 hash function [15].

Cryptographic Hash Functions serve a dual purpose: they can be used for data validation as well as user/device validation [16]. This is done by implementing digital signature schemes[17]. Without using Hash, the signature will be the same size as the message. Hash functions are widely used in backend databases not only for password management but also for index preparation, which speeds up record selection [18]. Table 1 shows currently available hashing algorithms.



Table 1: Existing Hashing Algorithms

Algorithm	Output size (bits)	Collision resistance	Preimage resistance	Speed
SipHash [19]	64-512	Strong	Strong	Fast
GOST [20]	256	Strong	Strong	Medium
BLAKE2 [21]	512	Strong	Strong	Fast
Whirlpool [22]	512	Strong	Strong	Medium
Tiger [23]	192	Strong	Strong	Fast
RIPEMD [24]	128-256	Weaker than SHA-256	Weaker than SHA-256	Medium
SHA-1 [25]	160	Weaker than SHA-256	Weaker than SHA-256	Fast
SHA-256 [26]	256	Stronger than SHA-1	Stronger than SHA-1	Fast
MD5 [27]	128	Weaker than SHA-1	Weaker than SHA-1	Fast

Data hiding and completeness are extremely critical. To address this issue, a combined DES and SHA-1 encryption system founded on the VC++ environment was created by Zhang et al [28]. The system offers a wide range of practicality since it uses the 3DES and RSA procedures for data encryption and the SHA-1 algorithm to confirm data integrity.

Ambedkar et al proposed an algorithm [29] each round used twenty-one steps to create a 64-bits fixed size hash code that is independent of any initial value. As a result, it is extremely difficult to reassemble the original message because no publicly known initial value was used. Logical operations are used and a three-bit circular left shift operation during hash function processes.

The paper is structured as follows: Section 2 outlines the proposed methodology, Section 3 details the experimental setup and results, and the final section contains the conclusion and suggestions for future work.

## 2. PROPOSED METHODOLOGY

In the current era of digitization, IoT has become the major factor to revolutionize the device automation with the use of internet. Use of IoT is spreading all over including, healthcare, home automation and in smart agriculture. These IoT devices can be controlled and operated by smart devices having low computation and memory resources. Use of sensor and actuators require high velocity of data transmission over the links. Data and commands need to be secured to avoid control, monitoring and operation of IoT devices by adversary.

In this research proposed algorithm, DES based Scalable Hashing Algorithm (DESSHA) offers highly secure cryptographic scheme to secure data transmission in terms of strong data integrity. DESSHA uses block processing mechanism of Data Encryption Standard (DES), a modern symmetric encryption algo. To maintain high security it uses expansion table, S-



Boxes and permutation table to deceive adversary while maintaining high security for the resource constrained IoT devices. DESSHA is able to produce secure hashes of input blocks of size as required by smart devices in any length to achieve scalability.



Figure 1: DESSHA Hashing

Figure 1 illustrates formal process of creating hash from the plaintext using DESSHA (hash function) to produce a fixed length hash (hexadecimal values). Input file is divided into blocks of 56 bits each, padded with zeroes if required.

Figure 2 illustrates the overall working of proposed framework for DESSHA, defined buffers  $g_1, g_2, g_3$  and  $g_4$  are initialized with particular hexadecimal values. Additionally, the corresponding  $g$  values are concatenated to form the buffers  $h_0$  and  $h_1$ . The inputs for the block processing function are  $h_0, h_1$ , and a key. Block processing performs all the standard operations of DES including several rounds of permutations, use of S-Boxes and series of XOR operations to make it undiscoverable using usual statistical calculations.

By using the left extractor '[' and right extractor ']', it separates the left and right 32 bits of the input block. Then, by conducting XOR operations with '[' and ']', respectively, the values of  $h_0$  and  $h_1$  are updated. Finally,  $g_1, g_2, g_3$  and  $g_4$  are concatenated to create the final encrypted text  $\mathcal{H}$ . Using specified buffers and block processing procedures, this procedure demonstrates the conversion of input data into encrypted blocks composed of hexadecimal hash values. Depending upon the register size of device left extractor '[' and right extractor ']' is able to produce variable length of hashes as cipher text. For a resource constrained smart device with register size 16 bits or 8 bits DESSHA will produce hash of 16 bits or 8 bits. While other state of the art algorithms like SHA family lacks the flexibility to produce variable length of hashes.

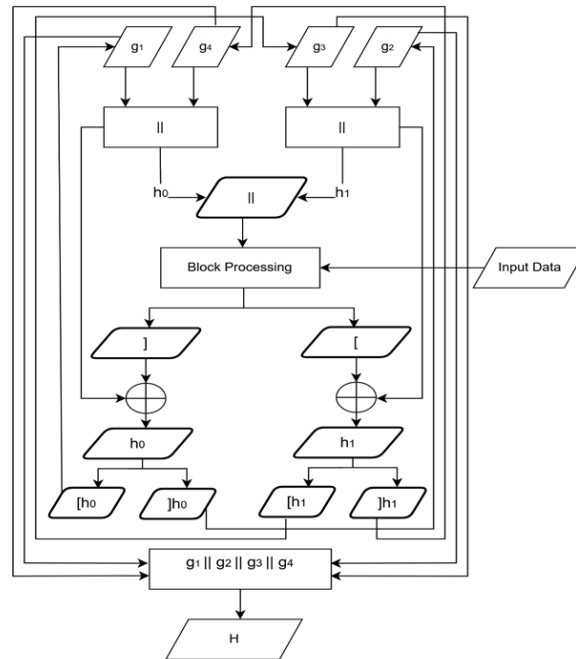


Figure 2: Proposed Framework for DESSHA

### Algorithm: DESSHA

We define stream concatenation and stream extraction.

- Stream Concatenation:

Let  $x$  and  $y$  be two streams of lengths  $m$  and  $n$  respectively where  $m, n \in \mathbb{Z}^+$ . Concatenation operator is denoted by  $||$  is defined as

$$x || y = x2^n + y$$

- Definition of Stream Left Extractor:

Let  $x$  be a bit stream  $x$  of length  $2k$ , where  $k \in \mathbb{Z}^+$ . Stream left extractor, denoted by  $[$  (left bracket), is defined as:

$$[x = \lfloor x \cdot 2^{-k} \rfloor$$

- Definition of Stream Right Extractor:

Let  $x$  be a bit stream  $x$  of length  $2k$ , where  $k \in \mathbb{Z}^+$ . Stream right extractor, denoted by  $]$  (right bracket), is defined as:

$$]x = x - 2^k \cdot \lfloor x$$



Let  $g_1, g_2, g_3$  and  $g_4$  be same some constants.

$$h_0 = g_3 || g_2 \qquad h_1 = g_1 || g_4$$

$$h = h_0 || h_2$$

let  $\mathcal{B}$  defined a block processing function and  $b$  be a block

$$\rho = \mathcal{B}(b)$$

where  $\rho$  is the processed block.

$$\rho_L = [\rho$$

$$\rho_R = ]\rho$$

We update  $g_1, g_2, g_3$  and  $g_4$  as follows

$$g_1 = [h_0$$

$$g_2 = ]h_0$$

$$g_3 = [h_1$$

$$g_4 = ]h_1$$

Hash digest  $\mathcal{H}$  is computed as

$$\mathcal{H} = g_1 || g_2 || g_3 || g_4$$

Initial values of  $g_i$  ( $i = 1,2,3,4$ ) are:

$$g_1 = x6A09$$

$$g_2 = xE667$$

$$g_3 = xBB67$$

$$g_4 = xEE85$$

### 3. RESULTS & DISCUSSION

There are a large number of resource constrained system which necessitate the strong cyber controls. Agriculture is considered as backbone of any country. IoT has reached out to facilitate the farmers to make it smart agriculture which is composed of several sensors and actuators. Initializing it from monitoring the crop facilitates to maintain several parameters including temperature, humidity, water level, nutrient levels (nitrogen, phosphorus, and potassium), and the conditions of irrigation components (fan, watering pump, and water



pump. Farmers will be able to increase production of crops with smart and automated decision making from the data being generated and processed by IoT components in smart agriculture.

Due to its vital role in intelligent decision-making data is of high importance to maintain high security to protect it against malfunctioning of equipment form adversary. Agriculture possesses paramount importance in society forming foundation of food security, this research takes initiative to secure data of smart agriculture system.

### 3.1 Dataset:

Cryptographic hash functions are used to maintain data integrity of data being sent over the medium. In this research DESSHA is applied to a dataset *IoT Agriculture 2024* [30]. Dataset is composed of several modern sensors and controller readings recorded in smart greenhouse at Tikrit University, Iraq. The dataset includes 13 different sensor readings over 37,923 rows for better management of IoT components. Table 2 shows minimal view of the dataset.

Table 2: Dataset - IoT Agriculture 2024

date	tempreat ure	humidit y	water _level	N			P			K			Fan_act		Watering		Water_p	
				N	P	K	OFF	ator_ON	ump_OFF	ump_ON	ump_act ump_OFF	ump_act ump_ON	ump_act ump_OFF	ump_act ump_ON				
08-02-24 6:10	41	63	100	255	255	255	0	1	1	0	1	0	1	0				
08-02-24 6:15	41	59	100	255	255	255	0	1	1	0	1	0	1	0				
08-02-24 6:20	41	62	100	255	255	255	0	1	1	0	1	0	1	0				
08-02-24 6:05	40	60	100	255	255	255	0	1	1	0	1	0	1	0				
08-02-24 6:00	39	61	100	255	255	255	0	1	1	0	1	0	1	0				
18-01-24 6:02	38	64	100	255	255	255	0	1	1	0	1	0	1	0				
18-01-24 6:07	38	65	100	255	255	255	0	1	1	0	1	0	1	0				
18-01-24 6:12	38	64	100	255	255	255	0	1	1	0	1	0	1	0				
18-01-24 5:02	37	62	100	255	255	255	0	1	1	0	1	0	1	0				
18-01-24 5:07	37	63	100	255	255	255	0	1	1	0	1	0	1	0				

### 3.2 Security Analysis

Cybersecurity of IoT devices is dependent upon the strong hashing algorithms. To assess the strength of hashing algorithm certain situations are imperative to be analyzed such as minimal changes in input to produce highly changed, random and unstructured cipher text. Following is some of the detailed criteria to evaluate effectiveness of data integrity of sensors data using DESSHA compared against theoretical values.

- Sensitivity of Hash

Sensitivity of hash is implementing single bit changes to an input message and supposed to produce altogether a highly dissimilar hash as the hash generated by original message. In our experiment DESSHA is applied to original message  $m_1$  to compute it hash as shown in figure 3 (a). To showcase DESSHA's sensitivity single bit changes are made in  $m_1$  in three different



cases: (b) flipping a single bit at random position to make it  $m_2$ ; (c) insertion of single bit at random position in bit string of original message to yield  $m_3$  and (d) deletion of single bit at random position of  $m_1$  to produce  $m_4$ . Table 3 clearly exhibits close alignment between mean values and variances of all three situation when compared with theoretical values achieving robustness and high sensitivity of DESSHA.

In addition, all 4 sections of figure 3 demonstrates 64-bit hashes of original and modified messages. The comparison highlights the altered bit positions in the hashes, marked with an asterisk (\*) at approximately 50% changed position of bits in hashes.

Table 3: Results of Sensitivity Testing

	Experimental Values			Theoretical Values		
	Bit Flipped	Bit Inserted	Bit Deleted	Bit Flipped	Bit Inserted	Bit Deleted
<b>mean</b>	0.4365	0.4981	0.5008	0.5	0.5	0.5
<b>variance</b>	0.00314	0.0043	0.0040	0	0	0

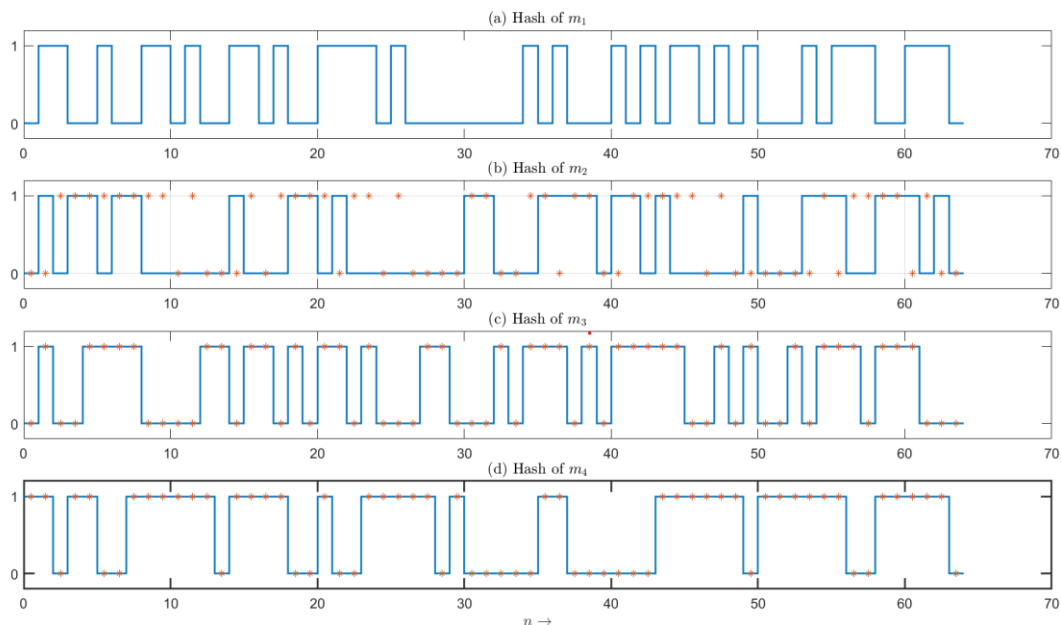


Figure 3: Sensitivity of Hashes

- Confusion Diffusion

In order to mislead the adversary, it is essential to distort the relationship between plaintext and ciphertext (confusion) and guaranteeing the slight changes in input to produce large



number of changes in ciphertext (diffusion). In table 4 DESSHA recurrently shows that mean of bits changed  $\bar{B}$  and the mean bit change probability  $P$  are very close to theoretical values. An important factor, confusion-diffusion measure  $I_{dc}$  presenting strength of DESSHA on closely comparable values of experimental output and theoretical results.

Table 4: Confusion Diffusion Analysis

	Experimental Values	Theoretical Values
Mean of bits changed ( $\bar{B}$ )	27.9330	32
Mean of bit change Probability ( $P$ )	0.4365	0.5
SD of bits changed ( $\Delta B$ )	11.3440	0
SD of bit change Probability ( $\Delta P$ )	0.1772	0
Confusion-Diffusion Measure ( $I_{dc}$ )	0.1204	0

- Uniform Distribution

Over the full bit stream of hash there are fair chances for every bit to be either 1 or 0, referred as uniform distribution. Attacker needs to process all possible options to break hash function which applies avalanche effect. In Figure 4, bit flip probabilities are very small in variance making it challenging to cryptanalyst as revealed in comparison of experimental and theoretical values.

Table 5: Uniform Distribution

	Experimental Values	Theoretical Values
Average Probability that a bit flips in hash ( $Q$ )	0.5005	0.5
$ Q - 0.5 $	0.0635	0
SD Probability that a bit flips in hash ( $\Delta Q$ )	0.0055	0

Smaller the  $|Q - 0.5|$  and  $\Delta Q$ , better is the avalanche effect.

Directing to table 5 the average probability of bit flips  $Q$  is nearly identical to the theoretical value of 0.5, which is proving uniform distribution. The minimal deviation  $|Q - 0.5|$  and low standard deviation of bit flip probability  $\Delta Q$  highlight the consistency and effectiveness of the system's avalanche effect.

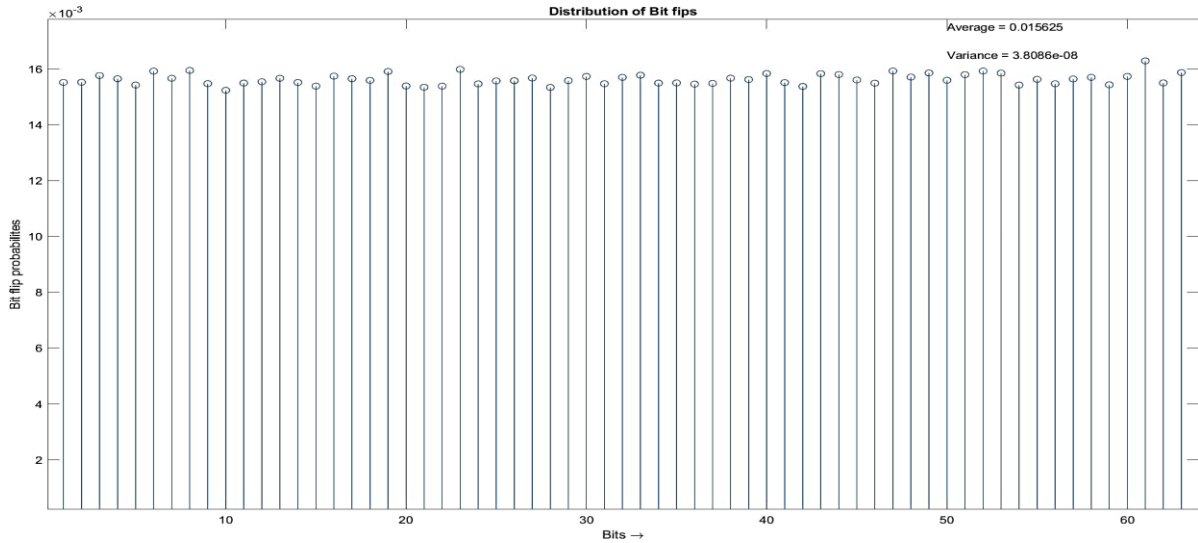


Figure 4: Uniform Distribution

- Collision Resistance

The major challenge in hash function is to guarantee computationally impractical to produce identical hash for any two distinct inputs. In figure 5, DESSHA's probability distribution is compared with theoretical predictions. There is very small difference between experimental and theoretical values which makes our algorithm highly resistant to collision. Consequently, proving it highly secure for the IoT devices to preserve data integrity.

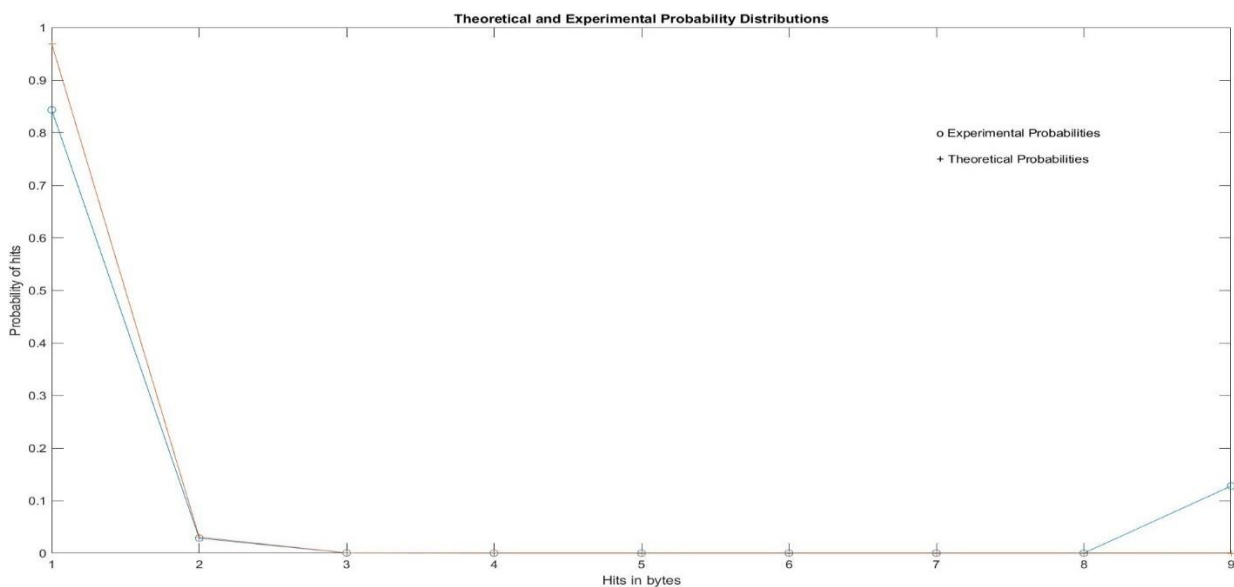


Figure 5: Collision Resistance



the uniform distribution property of a hash function could be evaluated using  $|Q-50\%|$

Bhattacharya distance is a statistical measure to quantify the similarity or dissimilarity between two probability distributions, expressed as  $D = (P^e || P^t)$  where  $P^e$  is probability distribution achieved on experimental setup whereas  $P^t$  is theoretical probability distribution. Smaller the value of  $D$  better the strict avalanche effect achieved that is  $D = 0.0684$ .

The absolute difference per byte between hash of original message  $\mathcal{H}$  and hash of modified message  $\overline{\mathcal{H}}$  can denoted as  $D_{byte}$ . Using DESSHA we obtained mean of  $D_{byte}$  as  $\overline{D}_{byte}^e = 74.5496$  which is very close to mean of theoretical value as  $\overline{D}_{byte}^t = 85.33$  that shows that DESSHA is capable of collision resistance.

### 3.3 Scalability

Let  $m$  represent the register size of any IoT device. Scalability enables the computation of hash values up to a size of  $64n \leq m$ , where  $n$  denotes the scalability factor. Our test results demonstrate that our proposed algorithm exhibits scalability, making it suitable for deployment across a wide range of smart devices in IoT. This scalability advantage is particularly noteworthy when compared to SHA-1 and SHA-256, which lack scalability. Figure 6 illustrates the comparison of DESSHA (Scalable), SHA-1 (Fixed), and SHA-256 (Fixed) hash sizes in the context of a 64-bit register. DESSHA demonstrates linear scalability, whereas SHA-1 and SHA-256 maintain constant hash sizes of 160 bits and 256 bits, respectively.

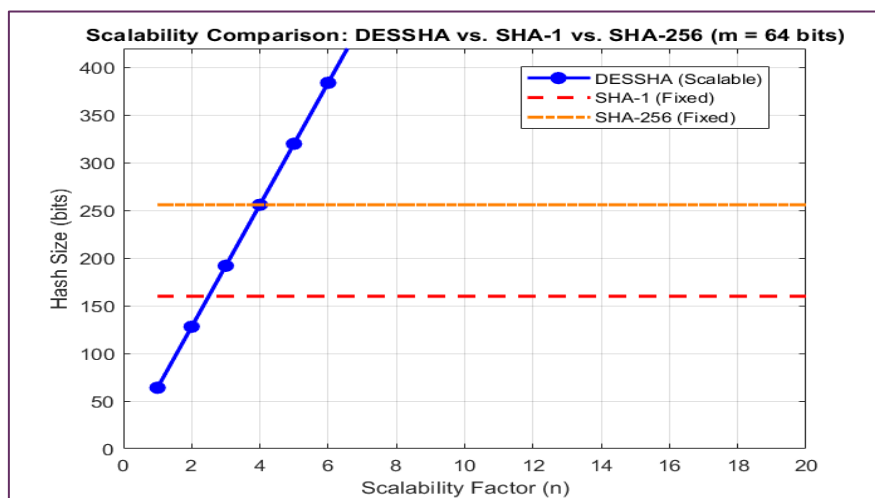


Figure 6: Visualization of Scalability in DESSHA vs SHA-1 vs SHA-256

## 4. Conclusion and Future Work

In the age digitalization and automation cybersecurity has become obligatory and challenging job. Data is being generated and sent at very large speed in IoT systems. Cryptographic hash



functions are one of the perfect security mechanisms to keep information safe. In this research a novel algorithm DESSHA is proposed which resistant to many statistical attacks and scalable to many smart devices. DESSHA is useful in message authentication for various types of smart devices, specifically for smart agriculture and E-commerce purposes. Our analysis presents that DESSHA has an ability to secure message with level of effectiveness closely aligned to theoretical approach in hash sensitivity, confusion diffusion, uniform distribution and collision resistance.

As the deployment of IoT systems are increasing everyday more security will be needed. DESSHA contributes its effectiveness on several type of IoT systems including smart home, healthcare as well as in industrial control systems in future.

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